

COMPUTATION OF LOCAL ISS LYAPUNOV FUNCTIONS WITH LOW GAINS VIA LINEAR PROGRAMMING

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ABSTRACT. In this paper, we present a numerical algorithm for computing ISS Lyapunov functions for continuous-time systems which are input-to-state stable (ISS) on compact subsets of the state space. The algorithm relies on a linear programming problem and computes a continuous piecewise affine ISS Lyapunov function on a simplicial grid covering the given compact set excluding a small neighborhood of the origin. The objective of the linear programming problem is to minimize the gain. We show that for every ISS system with a locally Lipschitz right-hand side our algorithm is in principle able to deliver an ISS Lyapunov function. For C^2 right-hand sides a more efficient algorithm is proposed.

1. Introduction. The concept of input-to-state stability (ISS) was first introduced by Sontag [20] in the late 1980s and has soon turned out to be one of the most influential concepts for characterizing robust stability of nonlinear systems. Several results on ISS can be found in [20, 21, 22]. In [24], different equivalent formulations of ISS are given. In particular, it is shown that the ISS property is equivalent to the existence of an ISS Lyapunov function. The ISS notion is also useful in stability analysis of large scale systems. Stability of large interconnected networks of systems can be analyzed by means of ISS small gain theorems if all the subsystems are ISS [6, 7, 8, 9, 16]. Motivated by these results, in this paper we focus on the computation of ISS Lyapunov functions for small systems, as the knowledge of ISS

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Lyapunov functions immediately leads to the knowledge of ISS gains which may be used in a small gain based stability analysis.

Based on [23, Lemma 2.10–2.14], it was shown in [4] that ISS Lyapunov functions in implication form may be calculated using a Zubov approach. An alternative Zubov type approach was developed in [17]. In these two papers the problem of computing ISS Lyapunov functions was transformed into the problem of computing robust Lyapunov functions for suitably designed auxiliary systems such as the auxiliary system described in Remark 3.2. This robust Lyapunov function can be characterized by the Zubov equation, a Hamilton-Jacobi-Bellman type partial differential equation, which can be solved numerically [3]. However, this approach only yields a numerical approximation of a Lyapunov function but not a true Lyapunov function. This means that the computed function is close to a Lyapunov function in some appropriate norm, but due to numerical errors it may fail to satisfy some of the properties a Lyapunov function is supposed to have, for instance it may not be decreasing along solution trajectories everywhere in its domain. For discrete time systems, following the same auxiliary system approach, true Lyapunov functions can be computed by a set oriented approach [11]. This numerical approach, however, does not carry over to the continuous time setting. Moreover, the detour via the auxiliary system introduces conservatism, since the resulting Lyapunov function and ISS gains strongly depend on the way the auxiliary system is constructed.

We thus propose a linear programming based algorithm for computing true ISS Lyapunov functions without introducing auxiliary systems. The approach to use linear programming for the computation of continuous, piecewise affine Lyapunov functions was first presented in [18]. In [12], it was proved that for exponentially stable equilibria the approach proposed in [18] always yields a solution. This result was extended to asymptotically stable systems [13], to asymptotically stable, arbitrarily switched, non-autonomous systems [14], and to asymptotically stable differential inclusions [2]. The approaches proposed in these papers yield true Lyapunov functions on compact subsets of the state space except possibly on arbitrarily small neighborhood of the asymptotically stable equilibrium. Mainly inspired by [2], in this paper we will propose an analogous linear programming based algorithm for computing true ISS Lyapunov functions for locally ISS systems.

The paper is organized as follows. In the ensuing Section 2, we introduce the notation and preliminaries. In Section 3 we present our algorithm along with a couple of auxiliary results needed in order to formulate the constraints in the resulting linear program. The algorithm is formulated in two variants for Lipschitz continuous and C^2 right hand sides. Section 4 contains the main results of the paper: we prove that upon successful termination the algorithm yields an ISS Lyapunov function outside a small neighborhood of the equilibrium, and that successful termination is guaranteed if the system admits a C^2 ISS Lyapunov function and the simplicial grid is chosen appropriately. In Section 5, we illustrate our algorithm and results by two numerical examples.

2. Notations and Preliminaries. Let $\mathbb{R}_+ := [0, +\infty)$ and for a set $\mathcal{C} \subset \mathbb{R}^n$ denote by $\text{cl}\mathcal{C}$ and $\text{int}\mathcal{C}$ its closure and interior respectively. For a vector $x \in \mathbb{R}^n$ we denote its transpose by x^\top . The standard inner product of $x, y \in \mathbb{R}^n$ is denoted by $\langle x, y \rangle$. We use the standard *norms* $\|x\|_p := (\sum_{i=1}^n |x_i|^p)^{1/p}$ for $p \geq 1$ and $\|x\|_\infty := \max_{i \in \{1, 2, \dots, n\}} |x_i|$ and let $B_p(z, r) := \{x \in \mathbb{R}^n \mid \|x - z\|_p < r\}$, $\text{cl}B_p(z, r) := \{x \in \mathbb{R}^n \mid \|x - z\|_p \leq r\}$ denote the open ball, closed ball of radius r

around z in the norm $\|\cdot\|_p$, respectively. The *induced matrix norm* is defined by $\|A\|_p := \max_{\|x\|_p=1} \|Ax\|_p$. By $\|u\|_{\infty,p} = \text{ess sup}_{t \geq 0} \|u(t)\|_p$ we denote the *essential supremum norm* of a measurable function $u : \mathbb{R}_+ \rightarrow \mathbb{R}^m$.

The *closed convex hull* of vectors $x_0, x_1, \dots, x_m \in \mathbb{R}^n$ is given by

$$\text{co}\{x_0, x_1, \dots, x_m\} := \left\{ \sum_{i=0}^m \lambda_i x_i \mid 0 \leq \lambda_i \leq 1, \sum_{i=0}^m \lambda_i = 1 \right\}.$$

A set of vectors $x_0, x_1, \dots, x_m \in \mathbb{R}^n$ is called *affinely independent* if $\sum_{i=1}^m \lambda_i (x_i - x_0) = 0$ implies $\lambda_i = 0$ for all $i = 1, 2, \dots, m$. This definition is independent of the numbering of the x_i , that is, of the choice of the reference point x_0 .

A continuous function $\alpha : \mathbb{R}_+ \rightarrow \mathbb{R}_+$ is said to be *positive definite* if it satisfies $\alpha(0) = 0$ and $\alpha(s) > 0$ for all $s > 0$. A positive definite function is of *class \mathcal{K}* if it is strictly increasing and of *class \mathcal{K}_∞* if it is of class \mathcal{K} and unbounded. A continuous function $\gamma : \mathbb{R}_+ \rightarrow \mathbb{R}_+$ is of *class \mathcal{L}* if $\gamma(r)$ is strictly decreasing to 0 as $r \rightarrow \infty$ and we call a continuous function $\beta : \mathbb{R}_+ \times \mathbb{R}_+ \rightarrow \mathbb{R}_+$ of *class \mathcal{KL}* if it is of class \mathcal{K}_∞ in the first argument and of class \mathcal{L} in the second argument.

In this paper we consider nonlinear perturbed systems described by the ordinary differential equation

$$\dot{x}(t) = f(x(t), u(t)) \quad (1)$$

with vector field $f : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}^n$, $f(0,0) = 0$, state $x(t) \in \mathbb{R}^n$, and perturbation input $u(t) \in \mathbb{R}^m$, $t \geq 0$. The admissible input values are given by $U_R := \text{cl}B_1(0, R) \subset \mathbb{R}^m$ for a constant $R > 0$ and the admissible input functions by $u \in \mathcal{U}_R := \{u \mid u : \mathbb{R}_+ \rightarrow U_R \text{ measurable}\}$. The solution corresponding to an initial condition $x(0) = x_0$ and an input $u \in \mathcal{U}_R$ is denoted by $x(\cdot, x_0, u)$.

For our algorithmic construction of Lyapunov functions, we need certain regularity properties of f which also determine certain inequalities imposed in the algorithm. To this end, we require one of the following two hypotheses.

(H1) The map $f : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}^n$ is locally Lipschitz continuous.

(H2) The vector field f is twice continuously differentiable.

Given a compact set $\mathcal{G} \subset \mathbb{R}^n$, as regards (H1) we fix the following notation: For each $u \in U_R$, $L_x(u)$ is a Lipschitz constant of the map $x \mapsto f(x, u)$, and for each $x \in \mathbb{R}^n$, $L_u(x)$ is a Lipschitz constant for the function $u \mapsto f(x, u)$. Moreover, by (H1) there exist constants \bar{L}_x and \bar{L}_u such that

$$\bar{L}_x \geq L_x(u) > 0, \bar{L}_u \geq L_u(x) > 0 \quad (2)$$

for all $x \in \mathcal{G}$, $u \in U_R$.

The following definition specifies the stability property we are considering in this paper.

Definition 2.1. The system (1) is called *locally input-to-state stable (ISS)*, if there exist $\rho_x > 0$, $\rho_u > 0$, $\gamma \in \mathcal{K}_\infty$ and $\beta \in \mathcal{KL}$ such that for all $\|x_0\|_2 \leq \rho_x$ and $\|u\|_\infty \leq \rho_u$

$$\|x(t, x_0, u)\|_2 \leq \beta(\|x_0\|_2, t) + \gamma(\|u\|_\infty), \quad \forall t \in \mathbb{R}_+. \quad (3)$$

If $\rho_x = \rho_u = \infty$, then the system (1) is called *input-to-state stable (ISS)*.

Observe that ISS implies that the origin is an equilibrium of (1) which is locally asymptotically stable for $u \equiv 0$. The function $\gamma \in \mathcal{K}_\infty$ is called an *ISS gain*.

It is known that the ISS property of (1) is equivalent to the existence of a smooth, i.e. C^∞ , ISS Lyapunov function for (1), see [23]. While this result guarantees

the existence of smooth ISS Lyapunov functions our numerical techniques will not generate a smooth function. In the following we will numerically construct piecewise affine and thus nonsmooth Lyapunov functions. In order to define these functions, we need a generalized notion of the gradient and for our purpose Clarke's subdifferential turns out to be useful. Since we are exclusively dealing with Lipschitz functions, we can use the following definition, cf. [5, Theorem 2.5.1].

Definition 2.2. For a Lipschitz continuous function $V : \mathbb{R}^n \rightarrow \mathbb{R}$, Clarke's subdifferential is given by

$$\partial_{Cl}V(x) := \text{co} \left\{ \lim_{i \rightarrow \infty} \nabla V(x_i) \mid x_i \rightarrow x, \nabla V(x_i) \text{ and } \lim_{i \rightarrow \infty} \nabla V(x_i) \text{ exist} \right\}. \quad (4)$$

Now we can state the definition of a nonsmooth ISS Lyapunov function.

Definition 2.3. Let $\mathcal{G} \subset \mathbb{R}^n$ with $\mathcal{G} \subset \text{clint } \mathcal{G}$ and $0 \in \text{int } \mathcal{G}$. Let $R > 0$. A Lipschitz continuous function $V : \mathcal{G} \rightarrow \mathbb{R}_+$ is said to be a (local) *nonsmooth ISS Lyapunov function* for system (1) on \mathcal{G} if there exist \mathcal{K}_∞ functions ψ_1, ψ_2, α and β such that

$$\psi_1(\|x\|_2) \leq V(x) \leq \psi_2(\|x\|_2) \quad (5)$$

$$\langle \xi, f(x, u) \rangle \leq -\alpha(\|x\|_2) + \beta(\|u\|_1) \quad (6)$$

hold for all $x \in \text{int } \mathcal{G}$, $u \in U_R$ and $\xi \in \partial_{Cl}V(x)$. If $\mathcal{G} = \mathbb{R}^n$ and $R = \infty$ then V is called a *global nonsmooth ISS Lyapunov function*. The function $\beta \in \mathcal{K}_\infty$ is called *Lyapunov ISS gain* or briefly *gain* in what follows. The gain is called *linear* if β is linear.

Remark 2.4. The particular norms chosen in the formulation of the ISS property in (3) or of an ISS Lyapunov function in (6) do not play a role from the conceptual point of view: as all norms in \mathbb{R}^n are equivalent, different norms will only lead to different numerical values of the gains. The particular formulations we have chosen will turn out to be useful in deriving easy estimates, see the last paragraph of the proof of Theorem 4.1.

In order to simplify the algorithm to be proposed in this paper, we will restrict ourselves to Lyapunov functions which satisfy (6) with linear functions $\alpha(s) = s$ and $\beta(s) = rs$ for some fixed $r > 0$. The following proposition shows that on compact subsets of the state space excluding a ball around the origin this can be done without loss of generality.

Proposition 2.5. *Let $\mathcal{G} \subset \mathbb{R}^n$ with $\mathcal{G} \subset \text{clint } \mathcal{G}$ and $0 \in \text{int } \mathcal{G}$ be compact. If there exists a Lipschitz continuous ISS Lyapunov function W for system (1) on \mathcal{G} , then for any $\epsilon > 0$ and $\sigma > 0$ there exist positive constants $C, r > 0$ such that $V(x) := CW(x)$ satisfies*

$$V(x) \geq \|x\|_2 \quad \forall x \in \mathcal{G} \setminus B_2(0, \epsilon) \quad (7)$$

and

$$\langle \xi, f(x, u) \rangle \leq -\sigma\|x\|_2 + r\|u\|_1 \quad \forall x \in \text{int } \mathcal{G} \setminus B_2(0, \epsilon), \xi \in \partial_{Cl}V(x), u \in U_R \quad (8)$$

with U_R from Definition 2.3.

Proof. From the assumption, there exist $\alpha, \beta \in \mathcal{K}_\infty$ such that W satisfies

$$\langle \xi, f(x, u) \rangle \leq -\alpha(\|x\|_2) + \beta(\|u\|_1), \quad \forall x \in \text{int } \mathcal{G} \setminus B_2(0, \epsilon), \xi \in \partial_{Cl}W(x), u \in U_R. \quad (9)$$

For the construction of C and r we now distinguish two cases.

Case 1: $\limsup_{s \rightarrow 0} \beta(s)/s$ is bounded. In this case we define

$$C = \min\{c \in \mathbb{R} \mid c\psi_1(\|x\|_2) \geq \|x\|_2 \text{ and } c\alpha(\|x\|_2) \geq \sigma\|x\|_2, \quad \forall x \in \mathcal{G} \setminus B_2(0, \epsilon)\}.$$

Then there exists a constant $r > 0$ satisfying

$$C\beta(\|u\|_1) \leq r\|u\|_1 \quad \text{for all } u \in U_R. \quad (10)$$

Case 2: $\limsup_{s \rightarrow 0} \beta(s)/s$ is unbounded. In this case we choose C as

$$C = \min\{c \in \mathbb{R} \mid c\psi_1(\|x\|_2) \geq \|x\|_2 \text{ and } c\alpha(\|x\|_2) \geq \sigma\|x\|_2 + \epsilon, \quad \forall x \in \mathcal{G} \setminus B_2(0, \epsilon)\}.$$

Then it is possible to find a constant $r > 0$ such that

$$C\beta(\|u\|_1) \leq r\|u\|_1 \quad (11)$$

for all $u \in U_R$ satisfying $C\beta(\|u\|_1) \geq \epsilon$.

In both cases, a straightforward calculation shows that $V(x) = CW(x)$ satisfies the desired inequalities. \square

Remark 2.6. It may not always be possible to choose $\epsilon = 0$ in Proposition 2.5. However, in general the linear programming approach to computing Lyapunov functions only works outside a neighborhood of the origin, anyway, cf. Remark 3.3, such that the need to remove $B_2(0, \epsilon)$ does not introduce additional limitations into our approach.

3. The algorithm. In this section we introduce an algorithm to compute a local ISS Lyapunov function defined on a compact set $\mathcal{G} \subset \mathbb{R}^n$ with $0 \in \text{int } \mathcal{G}$ and valid for perturbation inputs from $U_R \subset \mathbb{R}^m$. The algorithm uses linear programming and the representation of the function on a simplicial grid in order to obtain a numerical representation as a continuous, piecewise affine function. By taking into account interpolation errors, the algorithm yields a true Lyapunov function, not only an approximative one.

3.1. Definitions. We recall the following basic definitions: A simplex in \mathbb{R}^n is a set of the form $\Gamma = \text{co}\{x_0, x_1, \dots, x_j\}$, where x_0, x_1, \dots, x_j are affinely independent. The faces of Γ are given by $\text{co}\{x_{i_0}, x_{i_1}, \dots, x_{i_k}\}$, where $\{x_{i_0}, x_{i_1}, \dots, x_{i_k}\}$ ranges over the subsets of $\{x_0, x_1, \dots, x_j\}$. An n -simplex is generated by a set of $n + 1$ affine independent vertices. A collection \mathcal{S} of simplices in \mathbb{R}^n is called a simplicial complex, if

- (i) for every $\Gamma \in \mathcal{S}$, all faces of Γ are in \mathcal{S} ,
- (ii) for all $\Gamma_1, \Gamma_2 \in \mathcal{S}$ the intersection $\Gamma_1 \cap \Gamma_2$ is a face of both Γ_1 and Γ_2 (or empty).

Some authors consider the empty simplex to be a face of Γ , so that the last statement in (ii) is superfluous, but this will have no relevance in the present paper. The *diameter* of a simplex Γ is defined as $\text{diam}(\Gamma) := \max_{x, y \in \Gamma} \|x - y\|_2$.

We now return to our problem. We assume that $\mathcal{G} \subset \mathbb{R}^n$ may be partitioned into finitely many n -simplices $\mathcal{T} = \{\Gamma_\nu \mid \nu = 1, 2, \dots, N\}$, so that \mathcal{T} defines a simplicial complex. By assumption, we may also partition U_R into m -simplices $\mathcal{T}_u = \{\Gamma_\kappa^u \mid \kappa = 1, 2, \dots, N_u\}$ defining a simplicial complex. We briefly write $h_{x, \nu} = \text{diam}(\Gamma_\nu)$, $h_{u, \kappa} = \text{diam}(\Gamma_\kappa^u)$ and $h_x = \max_{\nu=1, 2, \dots, N} h_{x, \nu}$, $h_u = \max_{\kappa=1, 2, \dots, N_u} h_{u, \kappa}$. For each $x \in \mathcal{G}$ we define the *active index set* $I_{\mathcal{T}}(x) := \{\nu \in \{1, 2, \dots, N\} \mid x \in \Gamma_\nu\}$. For the simplices \mathcal{T}_u , we assume additionally that

$$\text{for each simplex } \Gamma_\kappa^u \in \mathcal{T}_u, \text{ the vertices of } \Gamma_\kappa^u \text{ are in the same closed orthant.} \quad (12)$$

Let $PL(\mathcal{T})$ denote the space of continuous functions $V : \mathcal{G} \rightarrow \mathbb{R}$ which are *affine* on each simplex, i.e., there are $a_\nu \in \mathbb{R}$, $w_\nu \in \mathbb{R}^n$, $\nu = 1, 2, \dots, N$, such that

$$V|_{\Gamma_\nu}(x) = \langle w_\nu, x \rangle + a_\nu \quad \forall x \in \Gamma_\nu, \Gamma_\nu \in \mathcal{T} \quad (13)$$

$$\nabla V_\nu := \nabla V|_{\text{int } \Gamma_\nu} = w_\nu \quad \forall \Gamma_\nu \in \mathcal{T}. \quad (14)$$

Let $\nabla V_{\nu,k}$ ($k = 1, 2, \dots, n$) denote the k -th component of the vector ∇V_ν for every $\Gamma_\nu \in \mathcal{T}$.

Similarly, we define $PL_u(\mathcal{T}_u)$. Observe that (12) implies that the map $u \mapsto \|u\|_1$ is contained in $PL_u(\mathcal{T}_u)$.

Remark 3.1. The algorithm will construct an ISS Lyapunov function $V \in PL(\mathcal{T})$. In particular, this means that the inequality (8) has to be satisfied. To this end, observe that from the definition it follows that for any function $V \in PL(\mathcal{T})$ Clarke's subdifferential is given by

$$\partial_{CI}V(x) = \text{co}\{\nabla V_\nu \mid \nu \in I_{\mathcal{T}}(x)\}, \quad x \in \text{int } \mathcal{G}. \quad (15)$$

Hence, for fixed $x \in \text{int } \mathcal{G}$ and $u \in U_R$ inequality (8) becomes

$$\langle \xi, f(x, u) \rangle \leq -\sigma\|x\|_2 + r\|u\|_1 \quad \forall \xi \in \text{co}\{\nabla V_\nu \mid \nu \in I_{\mathcal{T}}(x)\}.$$

Linearity of the scalar product in its first argument implies that this is equivalent to

$$\langle \nabla V_\nu, f(x, u) \rangle \leq -\sigma\|x\|_2 + r\|u\|_1 \quad \forall \nu \in I_{\mathcal{T}}(x). \quad (16)$$

An inequality of this type will be used for ensuring (8) in the algorithm.

Remark 3.2. With the help of suitable auxiliary functions η_1 and η_2 which are continuous and piecewise affine on each simplex, we may introduce the auxiliary system

$$\dot{x} = f_\eta(x, u) := f(x, u) - \eta_2(u)\eta_1(x). \quad (17)$$

Then, using arguments similar to [17], it can be shown that a robust Lyapunov function for (17) is an ISS Lyapunov function for (1). However, it turns out that for computational purposes this detour via the auxiliary system is not efficient, as it leads to an algorithm in which two linear programs have to be solved and furthermore introduces conservatism into the estimates. Therefore we will not explicitly use the auxiliary system. Our approach is based on the algorithm from [2] for the computation of robust Lyapunov functions. The way we adapt it to the ISS case is, however, inspired by the structure of (17).

3.2. Interpolation errors. As in [14, 2], the key idea for the numerical computation of a true Lyapunov function lies in incorporating estimates for the interpolation errors on \mathcal{T} — and in this paper also on \mathcal{T}_u — into the constraints of a linear program. In this section we analyze the error terms we need for this purpose.

Let $x \in \Gamma_\nu = \text{co}\{x_0, x_1, \dots, x_n\} \in \mathcal{T}$, $x = \sum_{i=0}^n \lambda_i x_i$, $1 \geq \lambda_i \geq 0$, $\sum_{i=0}^n \lambda_i = 1$ and $u \in \Gamma_\kappa^u = \text{co}\{u_0, u_1, \dots, u_m\} \in \mathcal{T}_u$, $u = \sum_{j=0}^m \mu_j u_j$, $1 \geq \mu_j \geq 0$, $\sum_{j=0}^m \mu_j = 1$.

The basic idea of the algorithm is to impose conditions on $V \in PL(\mathcal{T})$ in the vertices x_i of the simplices $\Gamma_\nu \in \mathcal{T}$ which ensure that the function V satisfies the inequalities (5) and (8) for $\sigma = 1$ on the whole set $\mathcal{G} \setminus B_2(0, \epsilon)$. Note that $V \in PL(\mathcal{T})$ is completely determined by its values at the vertices of the simplices in \mathcal{T} .

In order to ensure the properness condition (5), we impose the condition

$$V(x_i) \geq \|x_i\|_2, \quad (18)$$

for every vertex $x_i \in \Gamma_\nu$, $V(0) = 0$ and $V \in \mathcal{PL}(\mathcal{T})$. According to (18), for $x \in \Gamma_\nu$ we have

$$V(x) = \sum_{i=0}^n \lambda_i V(x_i) \geq \sum_{i=0}^n \lambda_i \|x_i\|_2 \geq \|x\|_2. \quad (19)$$

Note that (19) and $V(0) = 0$ can only be true if the origin is a vertex of the grid. This we always ensure in our computations.

In order to make sure that $V(x)$ satisfies (16) for all $x \in \Gamma_\nu \subset \mathcal{G}$, $u \in \Gamma_\kappa^u \subset U_R$ via imposing inequalities in the node values $V(x_i)$, we need to incorporate an estimate of the interpolation error into the inequalities. To this end, we demand that

$$\langle \nabla V_\nu, f(x_i, u_j) \rangle - r \|u_j\|_1 + \|\nabla V_\nu\|_1 A_{\nu, \kappa} \leq -\|x_i\|_2, \quad (20)$$

for all $i = 0, 1, \dots, n$, $j = 0, 1, \dots, m$. Here $A_{\nu, \kappa} \geq 0$ is a bound for the interpolation error of f in the points (x, u) with $x \in \Gamma_\nu \subset \mathcal{G}$, $u \in \Gamma_\kappa^u \subset U_R$, $x \neq x_i$, $u \neq u_j$.

Remark 3.3. Close to the origin the positive term $\|\nabla V_\nu\|_1 A_{\nu, \kappa}$ may become predominant on the left hand side of (20), thus rendering (20) infeasible. This is the reason for excluding a small ball $B_2(0, \epsilon)$ in the construction of V . Under certain conditions on f this problem can be circumvented by choosing suitably shaped simplices near the origin. In order to keep the presentation in this paper concise we do not go into details here and refer to [10], instead.

Since (20) will be incorporated as an inequality constraint in the linear optimization problem, we need to derive an estimate for $A_{\nu, \kappa}$ before we can formulate the algorithm. For this purpose we introduce the following Proposition 3.4. Here, for a function $g : \mathbb{R}^n \times \mathbb{R}^m \rightarrow \mathbb{R}$ which is twice continuously differentiable with respect to its first argument, we denote the Hessian of $g(x, u)$ with respect to x at z by

$$H_g(z, u) = \begin{bmatrix} \left. \frac{\partial^2 g(x, u)}{\partial x_1^2} \right|_{x=z} & \cdots & \left. \frac{\partial^2 g(x, u)}{\partial x_1 \partial x_n} \right|_{x=z} \\ \vdots & \ddots & \vdots \\ \left. \frac{\partial^2 g(x, u)}{\partial x_n \partial x_1} \right|_{x=z} & \cdots & \left. \frac{\partial^2 g(x, u)}{\partial x_n^2} \right|_{x=z} \end{bmatrix}.$$

For the first argument $x \in \Gamma_\nu$, let

$$H_x(u) := \max_{z \in \Gamma_\nu} \|H_g(z, u)\|_2, \quad (21)$$

and let $K_x : U_R \rightarrow \mathbb{R}_+$, \bar{K}_x , respectively, denote a bounded function and a positive constant satisfying

$$\max_{\substack{z \in \Gamma_\nu \\ r, s=1, 2, \dots, n}} \left| \frac{\partial^2 g(z, u)}{\partial x_r \partial x_s} \right| \leq K_x(u) \leq \bar{K}_x \quad (u \in U_R). \quad (22)$$

In the next proposition which is proved in a similar way to [2, Proposition 4.1, Lemma 4.2 and Corollary 4.3], we state properties of scalar functions $g : \mathcal{G} \times U_R \rightarrow \mathbb{R}$ or vector functions $g : \mathcal{G} \times U_R \rightarrow \mathbb{R}^p$ with respect to its first argument. Analogous properties hold with respect to the second argument.

Proposition 3.4. *Consider a convex combination $x = \sum_{i=0}^n \lambda_i x_i \in \Gamma_\nu$, $\Gamma_\nu = \text{co}\{x_0, x_1, \dots, x_n\}$, $\sum_{i=0}^n \lambda_i = 1$, $1 \geq \lambda_i \geq 0$, $u \in U_R$ and a function $g : \mathcal{G} \times U_R \rightarrow \mathbb{R}^p$ with components $g(x, u) = (g_1(x, u), g_2(x, u), \dots, g_p(x, u))$.*

(a) *If $g(x, u)$ is Lipschitz continuous in x with the bounds $L_x(u)$, \bar{L}_x from (2), then*

$$\left\| g \left(\sum_{i=0}^n \lambda_i x_i, u \right) - \sum_{i=0}^n \lambda_i g(x_i, u) \right\|_{\infty} \leq L_x(u) h_{x,\nu} \leq \bar{L}_x h_{x,\nu} \quad (23)$$

holds for all $x \in \mathcal{G}$, $u \in U_R$.

(b) If $g_j(x, u)$ is twice continuously differentiable with respect to x with the bound $H_x(u)$ from (21) on its second derivative for some $j = 1, 2, \dots, p$, then

$$\begin{aligned} \left| g_j \left(\sum_{i=0}^n \lambda_i x_i, u \right) - \sum_{i=0}^n \lambda_i g_j(x_i, u) \right| &\leq \\ \frac{1}{2} \sum_{i=0}^n \lambda_i H_x(u) \|x_i - x_0\|_2 \left(\max_{z \in \Gamma_\nu} \|z - x_0\|_2 + \|x_i - x_0\|_2 \right) &\leq H_x(u) h_{x,\nu}^2. \end{aligned} \quad (24)$$

Under the same differentiability assumption for all $j = 1, 2, \dots, p$, the estimate

$$\left\| g \left(\sum_{i=0}^n \lambda_i x_i, u \right) - \sum_{i=0}^n \lambda_i g(x_i, u) \right\|_{\infty} \leq nK_x(u) h_{x,\nu}^2 \leq n\bar{K}_x h_{x,\nu}^2 \quad (25)$$

holds for all $u \in U_R$ by assuming the bounds from (22).

(c) The analogous estimates hold for interpolation in u when the Hessian and the second derivatives are computed w.r.t. u . The corresponding constants will be denoted by $H_u(x)$, $L_u(x)$, $K_u(x)$ and \bar{K}_u .

Proof. We only prove the estimate (25) which is an immediate consequence of (24) and the estimate

$$H_x(u) = \max_{z \in \Gamma_\nu} \|H_g(z, u)\|_2 \leq nK_x(u) \leq n\bar{K}_x. \quad (26)$$

The proof of (26) follows from the following observation. Let $M \in \mathbb{R}^{n \times n}$, $|M|$ the matrix obtained by taking the absolute value component-wise, \bar{r} an upper bound for the absolute values of the entries in M and E the matrix with all entries equal to 1. Then we have $\|M\|_2 \leq \| |M| \|_2 \leq \bar{r} \|E\|_2 = n\bar{r}$. \square

3.3. The Algorithm. Now we have collected all the preliminaries to formulate the linear programming algorithm for computing an ISS Lyapunov function V for (1). In this algorithm, we introduce the values $V_{x_i} = V(x_i)$ as optimization variables. It is desirable to obtain an ISS Lyapunov function in which the influence of the perturbation represented by the value r in the gain $\beta(s) = rs$ in (8) is as small as possible. Since in general there is a tradeoff between σ and r in (8), see also [15], here we fix $\sigma = 1$. The objective of the linear program will then be to minimize the number r in (8) for $\sigma = 1$.

As explained in Remark 3.3, we only consider x satisfying $x \in \mathcal{G} \setminus B(0, \epsilon)$ for a small $\epsilon > 0$. To this end we define the subsets

$$\mathcal{T}^\epsilon := \{\Gamma_\nu \mid \Gamma_\nu \cap B(0, \epsilon) = \emptyset\} \subset \mathcal{T} \quad \text{and} \quad \mathcal{G}^\epsilon := \bigcup_{\Gamma_\nu \in \mathcal{T}^\epsilon} \Gamma_\nu. \quad (27)$$

Thus \mathcal{T}^ϵ is a triangulation of a compact set \mathcal{G}^ϵ away from the origin, with a distance of at least ϵ . We note that since $0 \in \text{int } \mathcal{G}$ we may choose $\epsilon > 0$ small enough so that the origin is separated from $\mathbb{R}^n \setminus \mathcal{G}$, loosely speaking $\mathcal{G} \setminus \mathcal{G}^\epsilon$ is inside of \mathcal{G} . In the remainder of the paper we will assume that $\epsilon > 0$ is chosen in this manner.

In the following algorithm, we will only impose the conditions (18) and (20) in those nodes x_i which belong to simplices $\Gamma \in \mathcal{T}^\epsilon$. Moreover, we use the estimates

of the interpolation errors $A_{\nu,\kappa}$ obtained from Proposition 3.4. Here we distinguish between vector fields f satisfying (H1) and (H2), respectively. While the stronger assumption (H2) allows for improved estimates, the weaker assumption (H1) applies to a larger class of vector fields f .

Algorithm

We solve the following linear optimization problem.

$$\text{Inputs: } \begin{cases} x_i, \|x_i\|_2 \text{ for all vertices } x_i \text{ of each simplex } \Gamma_\nu \in \mathcal{T}, \\ u_j, \|u_j\|_1 \text{ for all vertices } u_j \text{ of each simplex } \Gamma_\kappa^u \in \mathcal{T}_u, \\ h_{x,\nu} = \text{diam}(\Gamma_\nu) \text{ of each simplex } \Gamma_\nu \in \mathcal{T}^\epsilon, \\ h_{u,\kappa} = \text{diam}(\Gamma_\kappa^u) \text{ of each simplex } \Gamma_\kappa^u \in \mathcal{T}_u, \\ \text{choose } \bar{L}_x, \bar{L}_u \text{ from (2) if } f \text{ only satisfies (H1),} \\ \text{or choose } \bar{K}_x, \bar{K}_u \text{ from (25) and Proposition 3.4(c), respectively,} \\ \text{for } g(x, u) = f(x, u) \text{ from (1) if } f \text{ satisfies (H2).} \end{cases} \quad (28)$$

$$\text{Optimization variables: } \begin{cases} V_{x_i} \text{ for all vertices } x_i \text{ of each simplex } \Gamma_\nu \in \mathcal{T}, \\ C_{\nu,k} \text{ for } k = 1, 2, \dots, n \text{ and every } \Gamma_\nu \in \mathcal{T}^\epsilon, \\ r \in \mathbb{R}_+. \end{cases} \quad (29)$$

$$\text{Optimization problem:} \quad (30)$$

minimize r
subject to

- (A1) : $V_{x_i} \geq \|x_i\|_2$ for all vertices x_i
of each simplex $\Gamma_\nu \in \mathcal{T}$, and $V_0 = 0$.
(A2) : $|\nabla V_{\nu,k}| \leq C_{\nu,k}$ for each simplex $\Gamma_\nu \in \mathcal{T}^\epsilon$, $k = 1, 2, \dots, n$.
(A3) : $V_{x_i} < V_{x_j}$ for all vertices $x_i \in \partial(\mathcal{G} \setminus \mathcal{G}^\epsilon)$, $x_j \in \partial\mathcal{G}$.

For all vertices x_i of each simplex $\Gamma_\nu \in \mathcal{T}^\epsilon$, all vertices u_j
of each simplex $\Gamma_\kappa^u \in \mathcal{T}_u$, one of the following conditions
is required:

- (A4) : $\langle \nabla V_\nu, f(x_i, u_j) \rangle - r\|u_j\|_1 + A_{\nu,\kappa} \sum_{k=1}^n C_{\nu,k} \leq -\|x_i\|_2$,
where $A_{\nu,\kappa} = \bar{L}_x h_{x,\nu} + \bar{L}_u h_{u,\kappa}$ if f satisfies (H1) and
 $A_{\nu,\kappa} = n\bar{K}_x h_{x,\nu}^2 + m\bar{K}_u h_{u,\kappa}^2$ if f satisfies (H2).

Remark 3.5.

- (i) By (19), the condition (A1) yields $V(x) \geq \|x\|_2$ for $x \in \mathcal{G}^\epsilon$.
- (ii) The condition (A2) ensures that $V(x)$ does not explode and defines linear constraints on the optimization variables V_{x_i} , $C_{\nu,k}$.
- (iii) For (A3) note that $\partial(\mathcal{G} \setminus \mathcal{G}^\epsilon)$ is the inner boundary of \mathcal{G}^ϵ and $\partial\mathcal{G}$ is the outer boundary of \mathcal{G}^ϵ . Define M to be the maximum value of $V(x)$ at the inner boundary. The linear constraint (A3) implies that the sublevel set $\{x \in \mathcal{G} \mid V(x) \leq M\} \subset \text{int } \mathcal{G}$ can be assumed to include the set $B_2(0, \epsilon)$ as we may assume $V(x) < M$ in the interior of $\mathcal{G} \setminus \mathcal{G}^\epsilon$ without violating any constraints. If system (1) is locally ISS, then the condition (A3) is not necessary.

Remark 3.6. If the above linear optimization problem has a feasible solution, then the values $V_{x_i} = V(x_i)$ from this feasible solution at all vertices x_i of all simplices $\Gamma_\nu \in \mathcal{T}^\epsilon$ and the condition $V(x) \in PL(\mathcal{T}^\epsilon)$ uniquely define a continuous, piecewise affine function

$$V : \mathcal{G}^\epsilon \rightarrow \mathbb{R} \quad \text{on} \quad \mathcal{G}^\epsilon = \bigcup_{\Gamma_\nu \in \mathcal{T}^\epsilon} \Gamma_\nu. \quad (31)$$

Remark 3.7. It follows from Proposition 3.4 that instead of the term $n\bar{K}_x h_{x,\nu}^2 + m\bar{K}_u h_{u,\kappa}^2$ in (A4) when f fulfills (H2), one may use the sharper estimate

$$\begin{aligned} & \frac{n\bar{K}_x}{2} \left(\|x_i - x_0\|_2 \left(\max_{k=1,2,\dots,n} \|x_k - x_0\|_2 + \|x_i - x_0\|_2 \right) \right) \\ & + \frac{m\bar{K}_u(x_i)}{2} \left(\|u_j - u_0\|_2 \left(\max_{k=1,2,\dots,m} \|u_k - u_0\|_2 + \|u_j - u_0\|_2 \right) \right) \end{aligned}$$

with $K_u(x_i)$ from Proposition 3.4(c). The latter was used in our numerical experiments.

4. Main results. In this section we formulate and prove our two main results. We show that any feasible solution of our algorithm defines an ISS Lyapunov function on \mathcal{G}^ϵ and give conditions under which our algorithm will yield such a feasible solution. We start with the former.

Theorem 4.1. *If assumption (H1) or (H2) holds, and the linear optimization problem in (30) has a feasible solution, then the function V from (31) is an ISS Lyapunov function on \mathcal{G}^ϵ , i.e., it satisfies (5) and (6) for all $x \in \mathcal{G}^\epsilon$ and all $u \in U_R$.*

Proof. Consider convex combinations $x = \sum_{i=0}^n \lambda_i x_i \in \Gamma_\nu$, $\Gamma_\nu = \text{co}\{x_0, x_1, \dots, x_n\} \in \mathcal{T}^\epsilon$, $\sum_{i=0}^n \lambda_i = 1$, $1 \geq \lambda_i \geq 0$, and $u = \sum_{j=0}^m \mu_j u_j \in \Gamma_\kappa^u$, $\Gamma_\kappa^u = \text{co}\{u_0, u_1, \dots, u_m\} \in \mathcal{T}_u$, $\sum_{j=0}^m \mu_j = 1$, $1 \geq \mu_j \geq 0$.

First note that by (19) we have $V(x) \geq \|x\|_2$ for all $x \in \mathcal{G}^\epsilon$. Thus in (5) we may choose ψ_1 to be the identity and the existence of ψ_2 follows by Lipschitz continuity.

In order to prove inequality (8) for $\sigma = 1$ we compute

$$\begin{aligned} \langle \nabla V_\nu, f(x, u) \rangle &= \sum_{i=0}^n \lambda_i \langle \nabla V_\nu, f(x_i, \sum_{j=0}^m \mu_j u_j) \rangle \\ &+ \langle \nabla V_\nu, f(\sum_{i=0}^n \lambda_i x_i, \sum_{j=0}^m \mu_j u_j) \rangle - \sum_{i=0}^n \lambda_i \langle \nabla V_\nu, f(x_i, \sum_{j=0}^m \mu_j u_j) \rangle \\ &\leq \sum_{i=0}^n \lambda_i \langle \nabla V_\nu, f(x_i, \sum_{j=0}^m \mu_j u_j) \rangle + \|\nabla V_\nu\|_1 \left\| f(\sum_{i=0}^n \lambda_i x_i, u) - \sum_{i=0}^n \lambda_i f(x_i, u) \right\|_\infty \\ &\leq \sum_{i=0}^n \lambda_i \sum_{j=0}^m \mu_j \langle \nabla V_\nu, f(x_i, u_j) \rangle + \|\nabla V_\nu\|_1 \left\| f(\sum_{i=0}^n \lambda_i x_i, u) - \sum_{i=0}^n \lambda_i f(x_i, u) \right\|_\infty \\ &\quad + \sum_{i=0}^n \lambda_i \|\nabla V_\nu\|_1 \left\| f(x_i, u) - \sum_{j=0}^m \mu_j f(x_i, u_j) \right\|_\infty. \end{aligned}$$

According to Proposition 3.4, the constraints from (A4) ensure that V satisfies

$$\langle \nabla V_\nu, f(x, u) \rangle \leq - \sum_{i=0}^n \lambda_i \|x_i\|_2 + r \sum_{i=0}^m \mu_j \|u_j\|_1 \leq -\|x\|_2 + r\|u\|_1.$$

In the last step we have used the equality $\sum_{i=0}^m \mu_j \|u_j\|_1 = \|u\|_1$, which is true by assumption (12) and because we used the 1-norm for u . Indeed, this assumption ensures that the signs of the entries of u_j coincide in each entry location. Thus we have shown (8) with $\sigma = 1$ for all $x \in \mathcal{G}^\epsilon$ and all $u \in U_R$. \square

Now we turn to the second objective of this section. We derive conditions under which the linear programming problem has a feasible solution. To this end, we need a certain regularity property of the simplices in our grids. In order to formalize these, we need the following notation.

For each $\Gamma_\nu = \text{co}\{x_0, x_1, \dots, x_n\} \in \mathcal{T}^\epsilon$, let $y = x_0$, and define the $n \times n$ matrix $X_{\nu,y}$ by writing the components of the vectors $x_1 - y, x_2 - y, \dots, x_n - y$ as row vectors consecutively, i.e.,

$$X_{\nu,y} = (x_1 - y, x_2 - y, \dots, x_n - y)^\top. \quad (32)$$

Let $X_{\nu,y}^* := \|X_{\nu,y}^{-1}\|_2$. In Part (ii) of the proof of [2, Theorem 4.6] it is proved that $X_{\nu,y}^*$ is independent of the order of x_1, x_2, \dots, x_n . Moreover, $X_{\nu,y}^* = \lambda_{\min}^{-1}$ holds, where λ_{\min} is the smallest singular value of $X_{\nu,y}$. We define

$$X_\nu^* := \max_{y \text{ vertex of } \Gamma_\nu} \|X_{\nu,y}^{-1}\|_2, \quad \text{and} \quad X^* := \max_{\nu=1,2,\dots,N} X_\nu^*. \quad (33)$$

The regularity property now demands that we need to avoid grids with arbitrarily flat simplices. Formally, this means that there exists a positive constant $R_1 > 0$ such that all simplices $\Gamma_\nu \in \mathcal{T}^\epsilon$ in the considered grids satisfy the inequality

$$X_\nu^* \cdot \text{diam}(\Gamma_\nu) \leq X^* h_x \leq R_1, \quad (34)$$

for X_ν^* and X^* from (33), cf. [2, Remark 4.7], and $h_x = \max_{\nu=1,2,\dots,N} h_{x,\nu}$ from Section 3.1.

Theorem 4.2. *Consider a system (1) which satisfies (H1) or (H2) and which is ISS. Let $\epsilon > 0$ and $R_1 > 0$. Then, for all grids \mathcal{T}^ϵ and \mathcal{T}_u such that \mathcal{T}^ϵ satisfies (34) and for which h_x and h_u are sufficiently small, the linear programming problem from our algorithm has a feasible solution and delivers an ISS Lyapunov function $V \in PL(\mathcal{T}^\epsilon)$ on \mathcal{G}^ϵ .*

Proof. Since system (1) is ISS, there exists a C^2 ISS Lyapunov function $W : \mathcal{G} \rightarrow \mathbb{R}$, see [23]. Applying Proposition 2.5 we may without loss of generality assume that W satisfies (7) and (8) with $\sigma = 2$ and some $r > 0$.

Now consider an arbitrary but fixed $\Gamma_\nu = \text{co}\{x_0, x_1, \dots, x_n\} \in \mathcal{T}^\epsilon$. Let $y = x_0$, and define

$$W_{\nu,y} := \begin{pmatrix} W(x_1) & - & W(y) \\ W(x_2) & - & W(y) \\ & \vdots & \\ W(x_n) & - & W(y) \end{pmatrix}.$$

As in Part (iii) of the proof of [2, Theorem 4.6], we obtain

$$W_{\nu,y} - X_{\nu,y} \nabla W(y) := \frac{1}{2} \begin{pmatrix} \langle x_1 - y, H_W(z_1)(x_1 - y) \rangle \\ \langle x_2 - y, H_W(z_2)(x_2 - y) \rangle \\ \vdots \\ \langle x_n - y, H_W(z_n)(x_n - y) \rangle \end{pmatrix},$$

where $H_W(z)$ denotes the Hessian of $W(x)$ at z and $z_i = y + \xi_i(x_i - y)$ for some $\xi_i \in [0, 1]$.

Thus, using Proposition 3.4 (by ignoring the dependence on u), we get

$$\|W_{\nu,y} - X_{\nu,y} \nabla W(y)\|_2 \leq \frac{1}{2} n^{\frac{3}{2}} A h_x^2, \quad (35)$$

where $A := \max_{\substack{z \in \mathcal{G} \\ i,j=1,2,\dots,n}} \left| \frac{\partial^2 W}{\partial x_i \partial x_j}(z) \right|$. Applying Proposition 3.4 again, we obtain

$$\begin{aligned} \|X_{\nu,y}^{-1} W_{\nu,y} - \nabla W(x_i)\|_2 &\leq \|X_{\nu,y}^{-1} W_{\nu,y} - \nabla W(y)\|_2 + \|\nabla W(y) - \nabla W(x_i)\|_2 \\ &\leq \|X_{\nu,y}^{-1}\|_2 \|W_{\nu,y} - X_{\nu,y} \nabla W(y)\|_2 + \|\nabla W(y) - \nabla W(x_i)\|_2 \\ &\leq \|X_{\nu,y}^{-1}\|_2 \|W_{\nu,y} - X_{\nu,y} \nabla W(y)\|_2 + \max_{z \in \mathcal{G}} \|H_W(z)\|_2 h_x \\ &\leq n A h_x \left(\frac{1}{2} X_{\nu,y}^* n^{\frac{1}{2}} h_x + 1 \right). \end{aligned}$$

After these preliminary considerations, we now assign values to the variables V_{x_i} and $C_{\nu,k}$ of the linear programming problem from the algorithm and show that they fulfill the constraints.

For each vertex $x_i \in \Gamma_\nu \in \mathcal{T}^\epsilon$, we let $V(x_i) = V_{x_i} := W(x_i)$. Since W satisfies (7), it is obvious that $V(x_i) = V_{x_i} \geq \|x_i\|_2$ for $x \in \mathcal{T}^\epsilon$. It thus remains to show (A4) for some $r > 0$.

To this end, choosing one simplex $\Gamma_\nu = \text{co}\{x_0, x_1, \dots, x_n\} \in \mathcal{T}^\epsilon$ and letting $y = x_0$, we get

$$\nabla V_\nu = X_{\nu,y}^{-1} W_{\nu,y}, \quad (36)$$

since V is linear affine on the simplex Γ_ν and

$$V(x) = V(y) + \langle X_{\nu,y}^{-1} W_{\nu,y}, (x - y) \rangle = V(y) + W_{\nu,y}^\top (X_{\nu,y}^\top)^{-1} (x - y). \quad (37)$$

For the variables $C_{\nu,k}$, we set

$$C_{\nu,k} := \|\nabla V_\nu\|_2 = \|X_{\nu,y}^{-1} W_{\nu,y}\|_2, \quad k = 1, 2, \dots, n. \quad (38)$$

Thus $C_{\nu,k} \geq |\nabla V_{\nu,k}|$ for each $\Gamma_\nu \in \mathcal{T}^\epsilon$. Since $\nabla W(x)$ is bounded on \mathcal{G} and (34) holds, there exists a positive constant \bar{C} such that

$$\begin{aligned} C_{\nu,k} = \|X_{\nu,y}^{-1} W_{\nu,y}\|_2 &\leq \|X_{\nu,y}^{-1}\|_2 \max_{z \in \mathcal{G}_\epsilon} \|\nabla W(z)\|_2 h_x \\ &\leq R_1 \max_{z \in \mathcal{G}_\epsilon} \|\nabla W(z)\|_2 \doteq \bar{C} \end{aligned} \quad (39)$$

holds for all ν and k . From this analysis and the fact that W satisfies (8) with $\sigma = 2$, we obtain that

$$\begin{aligned} \langle \nabla V_\nu, f(x_i, u_j) \rangle - r \|u_j\|_1 &= \langle \nabla W(x_i) + \nabla V_\nu - \nabla W(x_i), f(x_i, u_j) \rangle - r \|u_j\|_1 \\ &\leq -2 \|x_i\|_2 + \|X_{\nu,y}^{-1} W_{\nu,y} - \nabla W(x_i)\|_2 \|f(x_i, u_j)\|_2 \\ &\leq -2 \|x_i\|_2 + n A h_x \left(\frac{1}{2} X_{\nu,y}^* n^{\frac{1}{2}} h_x + 1 \right) D, \end{aligned}$$

where $D := \max_{x \in \mathcal{G}, u \in U_R} \|f(x, u)\|_2 < \infty$.

Now let $h = \max\{h_x, h_u\}$. If (H1) holds, i.e., in the Lipschitz case, the linear constraint from (A4) is fulfilled whenever $h > 0$ is so small that for all vertices x_i of simplices in \mathcal{T}^ϵ

$$n A h \left(\frac{1}{2} X_{\nu,y}^* n^{\frac{1}{2}} h + 1 \right) D + n \bar{C} (\bar{L}_x + \bar{L}_u) h \leq \|x_i\|_2. \quad (40)$$

In case (H2), i.e., f is C^2 , the linear constraint from (A4) is satisfied if $h > 0$ is so small that

$$nAh\left(\frac{1}{2}X^*n^{\frac{1}{2}}h + 1\right)D + n\bar{C}(n\bar{K}_x + m\bar{K}_u)h^2 \leq \|x_i\|_2, \quad (41)$$

where \bar{K}_x, \bar{K}_u are the constants satisfying inequality (22) with respect to x, u respectively. Thus, the theorem is proved. \square

5. Examples. In this section we illustrate the algorithm by two examples. In order to highlight the fact that our algorithm minimizes the gain r in (8) for $\sigma = 1$, in our first example we compare the result of our algorithm with two piecewise affine Lyapunov functions, for which a closed-form expression is derived following the construction of the proof of Theorem 4.2. Our second example shows the result of our algorithm for an example for which no closed-form ISS Lyapunov function is known.

Example 5.1. We consider the following system which is adapted from [19]

$$\begin{aligned} \dot{x}_1 &= -x_1[1 - (x_1^2 + x_2^2)] + 0.1x_2u^2, \\ \dot{x}_2 &= -x_2[1 - (x_1^2 + x_2^2)], \end{aligned} \quad (42)$$

where $x \in \mathcal{G} = B_2(0, 0.588) \subset \mathbb{R}^2$, $u \in U_R = \{u \in \mathbb{R} : |u| \leq 4.41\}$.

For this example, we obtain two ISS Lyapunov functions V_1 and V_2 on \mathcal{G} based on two different functions $W_1(x), W_2(x)$ following the construction in the proof of Theorem 4.2 and compare them with the numerical ISS Lyapunov function V delivered by the algorithm.

- (1) For constructing a theoretical ISS Lyapunov function V_1 we start with the quadratic function candidate $W_1(x) = x_1^2 + x_2^2$. It is obvious that W_1 is twice differentiable. For this function we obtain for the dynamics from (42)

$$\begin{aligned} \langle \nabla W_1(x), f(x, u) \rangle &= 2x_1\dot{x}_1 + 2x_2\dot{x}_2 \\ &\leq -2\|x\|_2^2(1 - \|x\|_2^2) + 0.05u^2. \end{aligned} \quad (43)$$

Here $\alpha(\|x\|_2) = 2\|x\|_2^2(1 - \|x\|_2^2)$ is an increasing function whenever $\|x\|_2 \in [0, \frac{\sqrt{2}}{2}]$ and can thus be extended as a \mathcal{K}_∞ function for $\|x\|_2 > \frac{\sqrt{2}}{2}$. Hence, for $\beta(|u|) = 0.05|u|^2 \in \mathcal{K}_\infty$, W_1 is an ISS Lyapunov function for system (42) on \mathcal{G} .

Now we follow the proof of Theorem 4.2 in order to construct a piecewise affine Lyapunov function V_1 satisfying the constraints in our algorithm. To this end, let $\epsilon = 0.048$. Then the appropriate rescaling constant C in Proposition 2.5 is given by $C = C_1 = \frac{1}{\epsilon(1-\epsilon^2)}$. Indeed, replacing $W_1(x)$ by $V_1(x) := C_1W_1(x) = C_1(x_1^2 + x_2^2)$ we obtain

$$\begin{aligned} \langle \nabla V_1(x), \dot{x} \rangle &= 2C_1x_1\dot{x}_1 + 2C_1x_2\dot{x}_2 \\ &\leq -2C_1\|x\|_2^2(1 - \|x\|_2^2) + C_1\frac{1}{20}u^2 \\ &\leq -2\|x\|_2 + C_1\frac{1}{20}u^2, \end{aligned} \quad (44)$$

for $x \in \mathcal{G}^\epsilon$, and $|u| \leq 4.41$.

For u satisfying $|u| \leq 4.41$, we now need to find $r_1 > 0$ with

$$r_1|u| \geq \frac{1}{20}C_1u^2. \quad (45)$$

Since in the algorithm the objective is to minimize r , we select the minimal r satisfying this inequality which is given by $r_1 = \frac{4.41}{20\epsilon(1-\epsilon^2)} = 4.6044$.

Now, linear interpolation of this $W_1(x)$ on a sufficiently fine grid \mathcal{T} yields the desired function $V_1(x) = C_1W_1(x)$, which is plotted in Figure 4.

- (2) Since in the construction in the proof of Theorem 4.2 we rescale the function via Proposition 2.5 to satisfy $W(x) \geq \|x\|_2$, it appears reasonable to start with $W_2(x) = \|x\|_2$ as a Lyapunov function candidate. Following the same steps as in (1), we can show that W_2 is also an ISS Lyapunov function. A rescaling $V_2(x) := C_2W_2(x)$ along Proposition 2.5 yields $C_2 = \frac{2}{1-\epsilon^2}$ and $r_2 = \frac{4.41}{10(1-\epsilon^2)} = 0.4410$. The resulting interpolated V_2 is shown in Figure 5.
- (3) From the algorithm, we get the numerical ISS Lyapunov function V shown in Figure 6 with $r = 0.420909$.

The simplicial complex is obtained in the same way as in [10, Section 2]: Let N_1, N_2, N_3, N_4 be positive integers such that $\Omega = [-N_1, N_2] \times [-N_3, N_4]$.

First, we pick points $x = (x_1, x_2) \in \Omega \cap \mathbb{Z}^2$ as vertices. The resulting triangulation of Ω is shown in Figure 1.

Second, let M_1, M_2, M_3, M_4 be positive integers, typically much smaller than N_1, N_2, N_3, N_4 , and $\Omega_1 = [-M_1, M_2] \times [-M_3, M_4] \subset \Omega$. On the small neighborhood Ω_1 around the equilibrium, the previous triangulation is replaced by triangles with one vertex in the origin and two vertices on the boundary of Ω_1 . The resulting second triangulation of Ω for $M_1 = \dots = M_4 = 2$ is shown in Figure 2.

Third, by the map $F : \mathbb{R}^2 \rightarrow \mathbb{R}^2$ with

$$F(x) = \begin{cases} \rho x \|x\|_\infty^2 / \|x\|_2 & \text{for } x \neq 0, \\ 0 & \text{for } x = 0 \end{cases} \quad (46)$$

we transfer the vertices from the second triangulation into new vertices from which we construct the simplices used in computation, cf. Figure 3. Due to the spherical shape of this triangulation, for $x \in B_2(0, \epsilon)$ we have $V(x) \leq \max_{x \in \partial(\mathcal{G} \setminus \mathcal{G}^\epsilon)} V(x)$. Thus the set $B_2(0, \epsilon)$ is a subset of the level set $\{x \in \mathcal{G} \mid V(x) \leq \max_{x \in \partial(\mathcal{G} \setminus \mathcal{G}^\epsilon)} V(x)\}$. The resulting set \mathcal{G} is shown in Figure 3. In the map F , the parameter $\rho > 0$ controls the size of the resulting vertices. For our computations we used $N_i = 7$, $M_i = 2$, $i = 1, 2, 3, 4$, and $\rho = 0.012$.

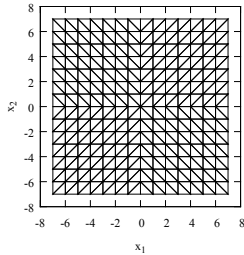


FIGURE 1. The first triangulation of Ω .

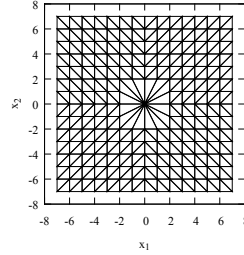


FIGURE 2. The second triangulation of Ω , $M_1, M_2, M_3, M_4 = 2$.

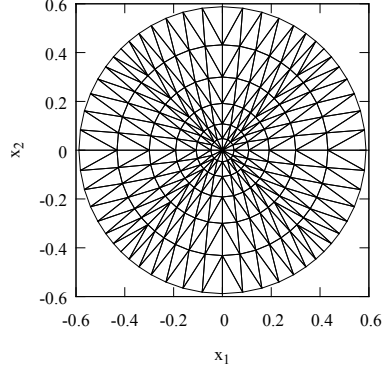


FIGURE 3. Simplices used in the computation.

The triangulation of U_R is obtained by the same method in 1d with $N_1 = N_2 = 21$, $M_1 = M_2 = 0$ and using the map $G : \mathbb{R} \rightarrow \mathbb{R}$ with

$$G(u) = \gamma u|u| \quad (47)$$

instead of F and $\gamma = 0.01$.

As expected, the optimization based algorithmic approach yields the smallest possible gain parameter r . For finer grids, even smaller values of r can be obtained and it appears that r converges to a lower bound $\underline{r} > 0.4$. However, we are not aware of an analytical method to confirm this numerically computed lower bound.

In Figures 8–9 we include a comparison of level curves of the calculated ISS Lyapunov function V with these of the theoretical obtained functions V_1 and V_2 . Note that the ISS Lyapunov function is not unique, but the calculated one is more similar to V_2 .

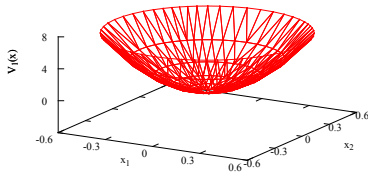


FIGURE 4. Theoretical ISS Lyapunov function $V_1(x)$ based on $W_1(x) = x_1^2 + x_2^2$ for system (42), $\epsilon = 0.048$, $r_1 = 4.6044$.

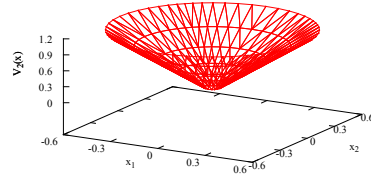


FIGURE 5. Theoretical ISS Lyapunov function $V_2(x)$ based on $W_2(x) = \|x\|_2$ for system (42), $\epsilon = 0.048$, $r_2 = 0.4410$.

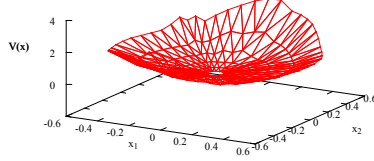


FIGURE 6. Numerical ISS Lyapunov function $V(x)$ delivered by the algorithm for system (42), $\epsilon = 0.048$, $r = 0.420909$.

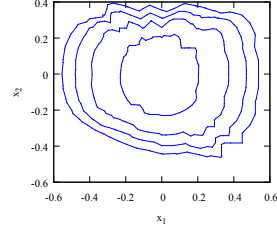


FIGURE 7. Level curves of ISS Lyapunov function V for system (42) at values 0.2732, 0.5464, 0.8196, 1.0928.

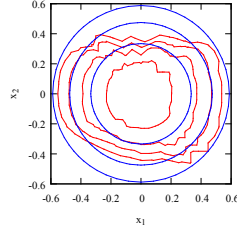


FIGURE 8. Level curves of ISS Lyapunov function V for system (42) at values 0.2732, 0.5464, 0.8196, 1.0928 and level curves of ISS Lyapunov function V_1 at values 2.40654, 4.81308, 7.21963 (circles).

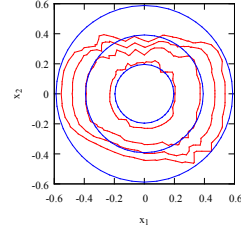


FIGURE 9. Level curves of ISS Lyapunov function V for system (42) at values 0.2732, 0.5464, 0.8196, 1.0928 and level curves of ISS Lyapunov function V_2 at values 0.3929067, 0.7858133, 1.17872 (circles).

Example 5.2 (Synchronous generator with varying damping). We consider the following model adapted from [1] which is described by

$$\begin{aligned} \dot{x}_1 &= x_2, \\ \dot{x}_2 &= -x_2 - \sin(x_1 + u) + \sin(u), \end{aligned} \quad (48)$$

on $\mathcal{G} = B_2(0, 2.352) \in \mathbb{R}^2$, $U_R = [-0.3, 0.3]$. For the triangulation of \mathcal{G} , we let $N_i = 14$ ($i = 1, 2, 3, 4$), $M_i = 1$ and utilize the map from (46) with $\rho = 0.012$. The triangulation of U_R is obtained with $N_1 = N_2 = 5$, $M_1 = M_2 = 0$, and the map from (47) with $\gamma = 0.012$. By the algorithm, we get the numerical ISS Lyapunov function V shown in Figure 10 for system (48). Some level curves of the ISS Lyapunov function V are shown in Figure 11. Note that for this example an analytical ISS Lyapunov function is not known and that our numerical analysis yields a numerical value for the (in our approach) linear ISS gain.

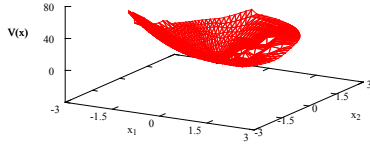


FIGURE 10. Local ISS Lyapunov function $V(x)$ given by the algorithm for system (48), $\epsilon = 0.012$, $r = 20.4783$.

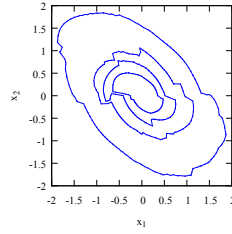


FIGURE 11. Level curves of ISS Lyapunov function V for system (48) at values 4.208675, 8.41735, 12.626025, 16.8347.

Remark 5.3. In our algorithm, we construct grids on U_R and \mathcal{G} , respectively. If \mathcal{G} is a two-dimensional set and the number of vertices for gridding U_R increases by 1, then the number of constraints in the linear program increases at least by $(48 + 6(N_v + 2)(N_v - 4) + 12(N_v - 2))$, where $N_v = \min\{N_{v,1}, N_{v,2}\}$ and $N_{v,i}$ is the number of vertices intersecting the x_i -axis. Similarly, the number of constraints increases in higher space dimensions. Thus, the gridding of U_R renders the number of constraints much larger than the number of optimization variables. It is hence much faster to solve the corresponding dual optimization problem than to solve the primal problem. For numerical computations we used the GNU Linear Programming Kit (GLPK)¹, Gurobi² and CPLEX³ respectively. We experienced that Gurobi and CPLEX carry out a significantly better preprocessing of the constraints which eliminates much more redundant ones and thus both methods can solve the optimization problem much faster than GLPK.

Since $f(x, u)$ and the interpolation errors on \mathcal{T} and \mathcal{T}_u are incorporated in constraints (A4), the grids of x, u have an influence in computing the gain parameter. From numerical experience, the grid of x plays a more important role in obtaining a good estimate of r via the solution of the linear optimization problem (30).

6. Conclusions. In this paper, we proposed a new method of computing a continuous piecewise affine ISS Lyapunov function for a dynamic system with input perturbation (1). For suitable triangulations of the state space and the input perturbation space, the algorithm delivers a true ISS Lyapunov function with a gain function for system (1) on a compact subset of the state space minus a small neighborhood of the origin (Theorem 4.1). Such an ISS Lyapunov function satisfies a linear inequality. We think the new computational approach will be helpful in analyzing stability of interconnected systems which are locally ISS, which is one topic of our future research. If system (1) has a C^2 ISS Lyapunov function, then there exist triangulations such that the algorithm has a feasible solution (Theorem 4.2). It is known that if system (1) is ISS, there exists a smooth ISS Lyapunov function [23]. Therefore, our proposed algorithm always has a feasible solution.

¹<http://www.gnu.org/software/glpk/>

²<http://www.gurobi.com/>

³<http://www.ibm.com/software/commerce/optimization/cplex-optimizer/>

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